

Computer organisation and Architecture.

Summary:-

- Digital components
 - ALU
 - Memory
 - Registers
- Instruction set Architecture
 - Format
 - Encoding
 - Addressing mode.
- Machine & Assembly language programming

Computer organisation:-

Question:- How to organise all the digital components such that a program could be executed?

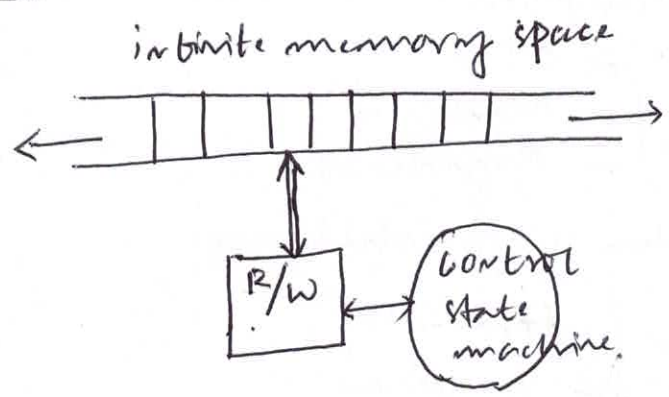
The principle:- stored program approach

- The idea:-
- ① Store the program in memory
(Set of well defined instructions) → unambiguous
 - ② Fetch the instruction (~~one by one~~)
 - ③ Decode the instruction to identify control signals and operand.
 - ④ Fetch the operand
 - ⑤ Execute the instruction using control signals and operand
 - ⑥ Store back the results
(Changing the current-state of the machine)

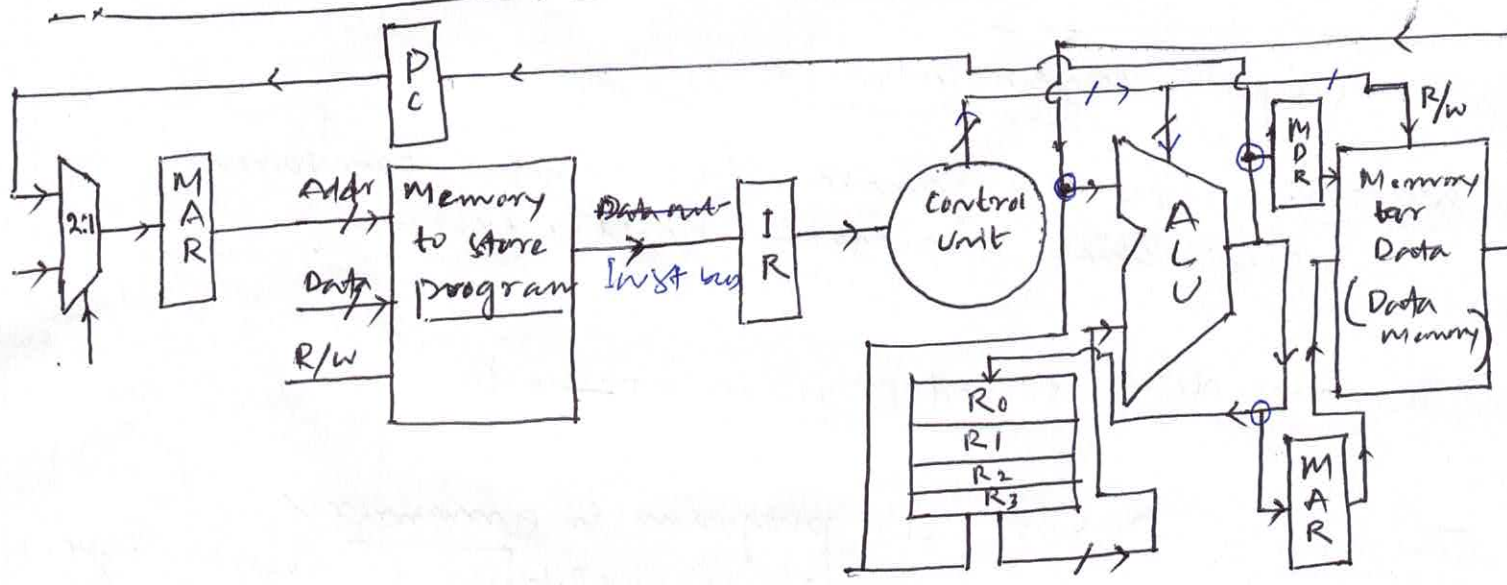
②

Von-Neumann Architecture | Turing Machine -

Turing Machine :-



Von-Neuman Architecture :-



- The connections in the above figure are conceptual
- The exact data & control paths with sufficient steering logic & interconnect will be drawn later on.

The Microoperations :-

- The micro steps required to execute an instruction through the five major steps.

Micro operations for ^{"HACK" (Hard to Tet's processor)} over processor :-

(3)

① Storing program at the first place

- This is a work of resource manager (operating system)

MAR \leftarrow Address of the stored program

Memory address bus \leftarrow [MAR]

Memory data bus \leftarrow ~~Addr~~ Instruction as a byte or word

All these microoperations be repeated to load a program from hard disk to main memory

(in the modern ^{"processor"} computer this is being done by DMA - Direct memory Access)

* In the HACK processor the program is being loaded ^{Access} into RAM memory through the means of simulator

② Instruction Fetch :-

(IF)

- ~~Address of that instruction is required~~

- ~~place to keep the instruction once fetched~~

- ~~Any post processing.~~

- Requirement :-

- Address of the instruction to be fetched.

- place to keep the fetched instruction.

- post processing (updating PC content)

[Address bus] \leftarrow Content of [PC]

~~REP~~ R/W \leftarrow 0

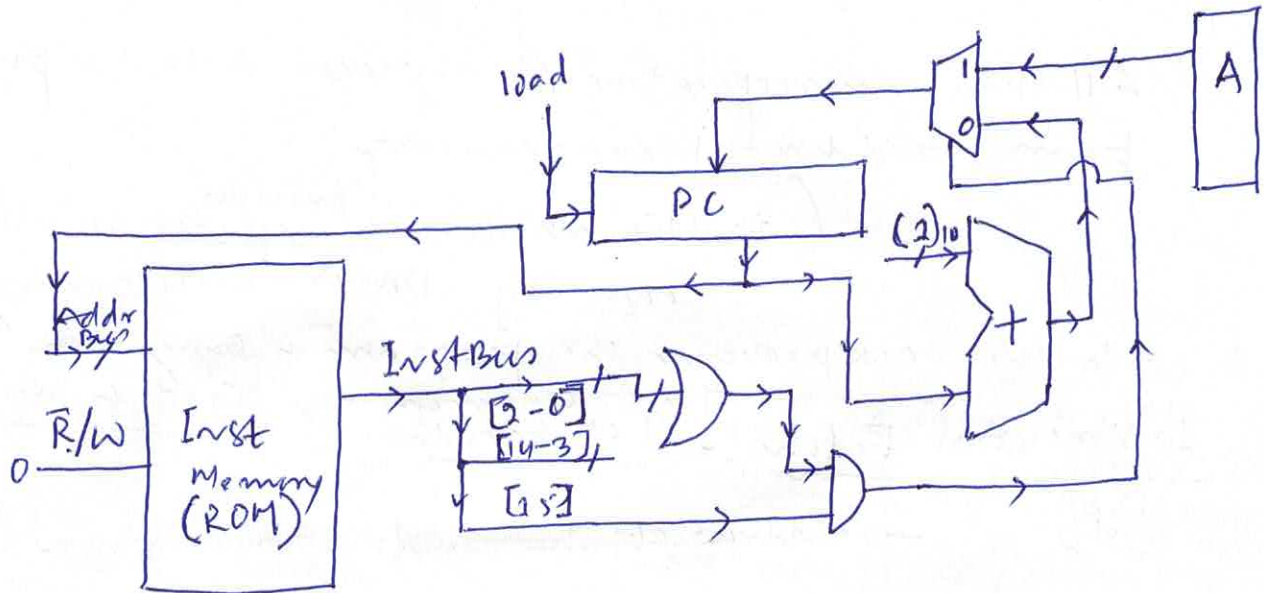
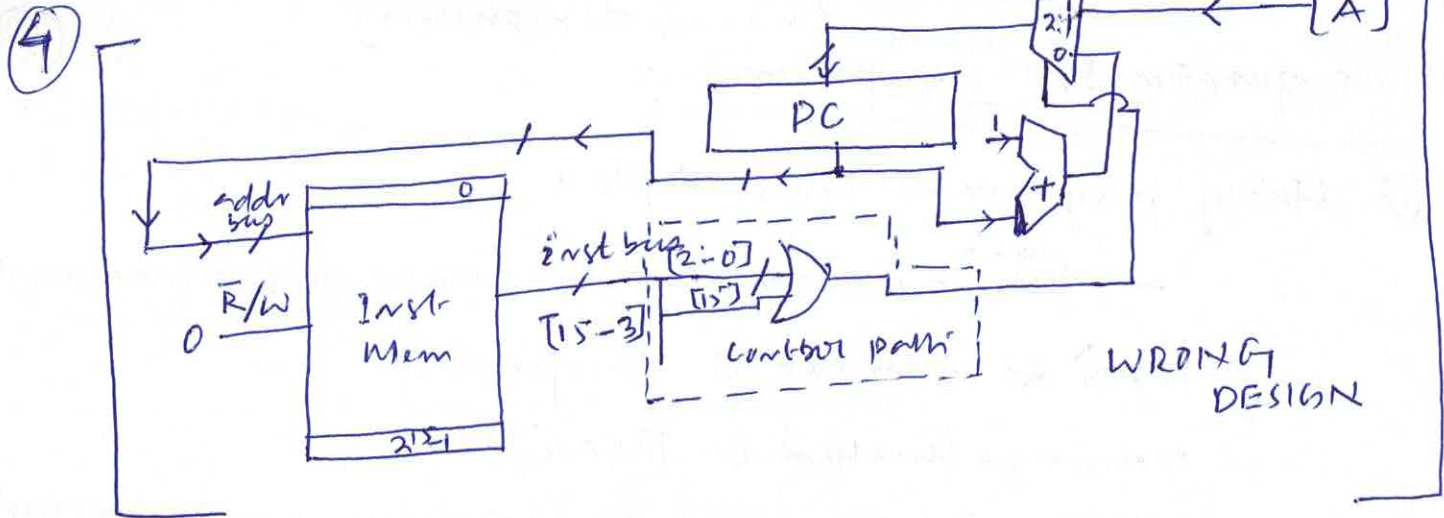
[Instruction bus] \leftarrow Memory [Add bus]

if instruction [21|0] \neq 0 then

PC \leftarrow [A]

else

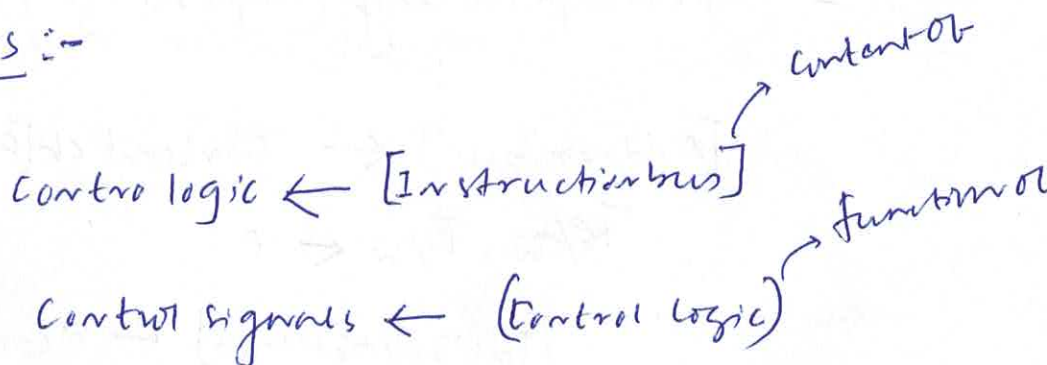
PC \leftarrow PC + 1



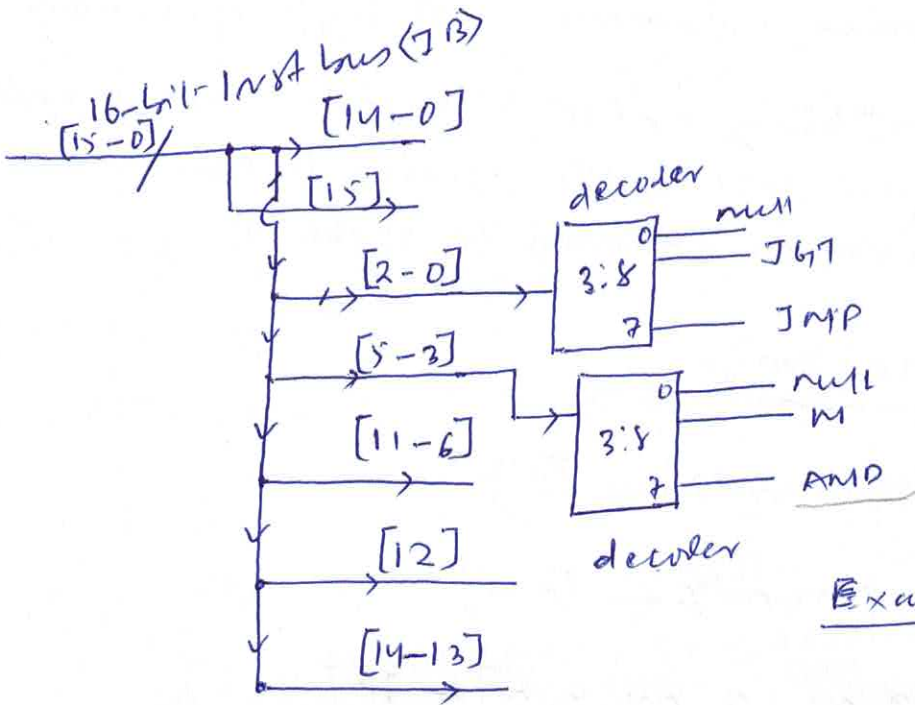
Instruction Decode: -(10)

Decoding is essential to understand the control signals and operands.

Micro operations :-



"Partial" control path design :-



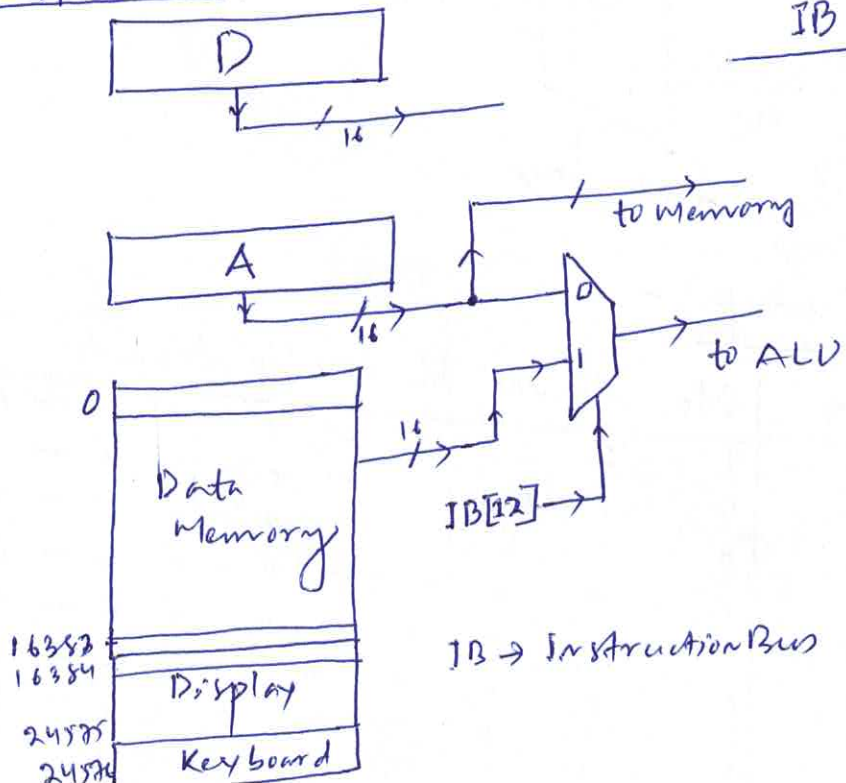
Example:-
 x86 (CISC) - microprogram control
 RISC - Hardware control

Operand Fetch (OF)

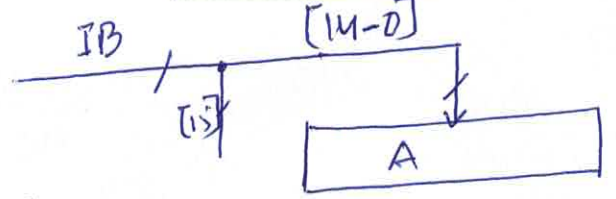
Two register → D - data register
 A - can be used as data register
 also as address register

one memory → M - memory

ALU instruction:-



Address/Mem instruction



- 16352
- 16354
- 24575
- 24576
- Display
- Keyboard

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Execute (EX)

- To perform arithmetic and logic operation
- The necessary control signals are appropriately connected to steer the operand to ALU and execute the instruction according to op-code.

micro operations :-

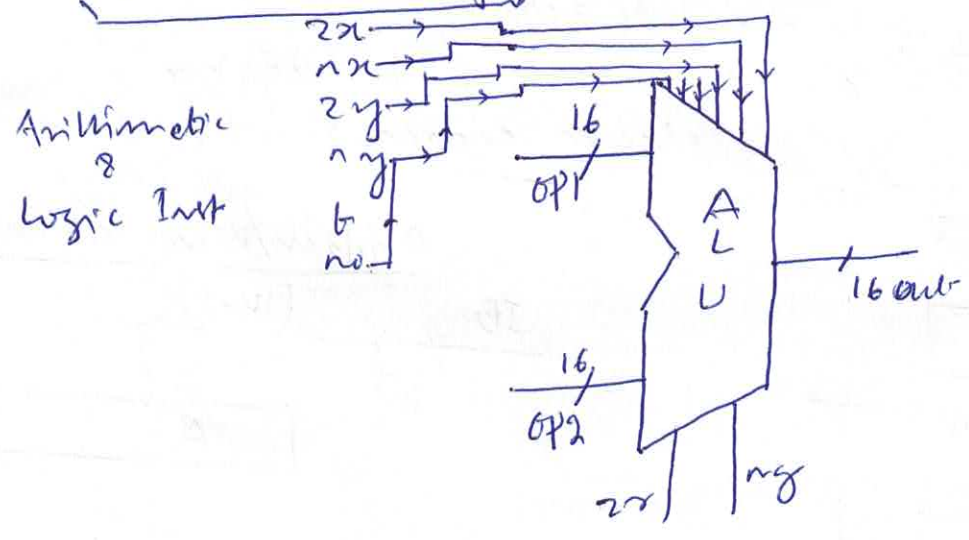
$$\text{operand}_1 \leftarrow [D]$$

$$\text{operand}_2 \leftarrow [AM]$$

$$[zx, nx, zy, ny, b, no] \leftarrow IB[11-6]$$

$$\text{data bus} \leftarrow ALU \text{ out}$$

Data path Design :-



Data path for Jump :-

J_1	J_2	J_3	J_4	J_5	J_6	J_7	
1	0	0	0	0	0	0	— null
0	0	0	1	0	0	0	— out > 0
0	0	1	0	0	0	0	— out = 0
0	1	0	0	1	0	0	— out >= 0
0	1	1	0	0	1	0	— out < 0
1	0	0	0	0	0	1	— out ≠ 0
1	0	1	0	0	0	0	— out ≤ 0
1	1	0	0	0	0	0	— JMP

$zr = 1$ if $out = 0$ else $zr = 0$
 $ng = 1$ if $out < 0$ else $ng = 0$

$out = 0 \rightarrow zr = 1$

$out < 0 \rightarrow ng = 1$

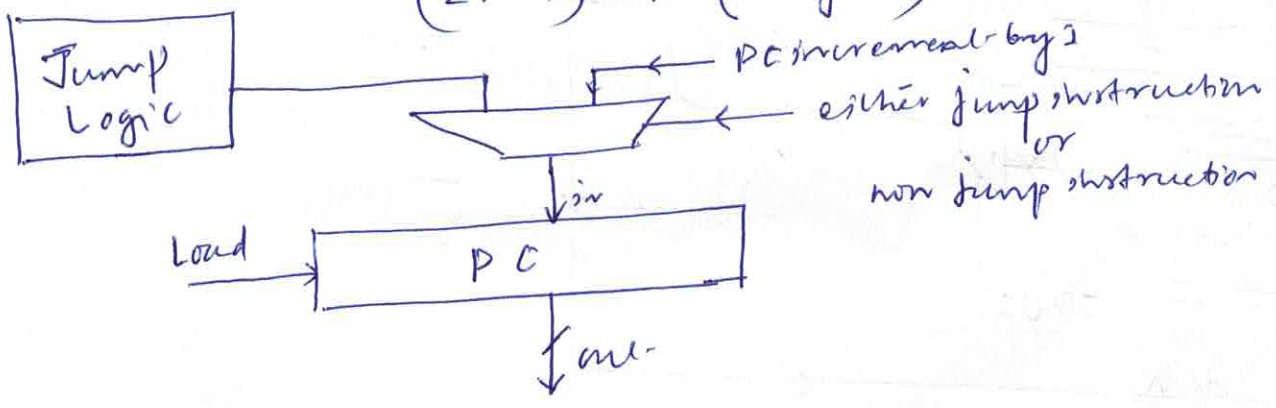
$out \neq 0 \Rightarrow out \neq 0 \ \& \ out \neq 0$
 $zr = 0 \ \& \ ng = 0$

$out \neq 0 \Rightarrow out = 0 \ \text{OR} \ out > 0$
 $(zr = 1) \ \text{OR} \ (zr = 0 \ \& \ ng = 0)$

$out \neq 0 \rightarrow zr = 0$

$out \leq 0 \Rightarrow out = 0 \ \text{OR} \ out < 0$

$(zr = 1) \ \text{OR} \ (ng = 1)$



Home work :- Complete the design of JUMP Logic
 using the above conditional control signals.

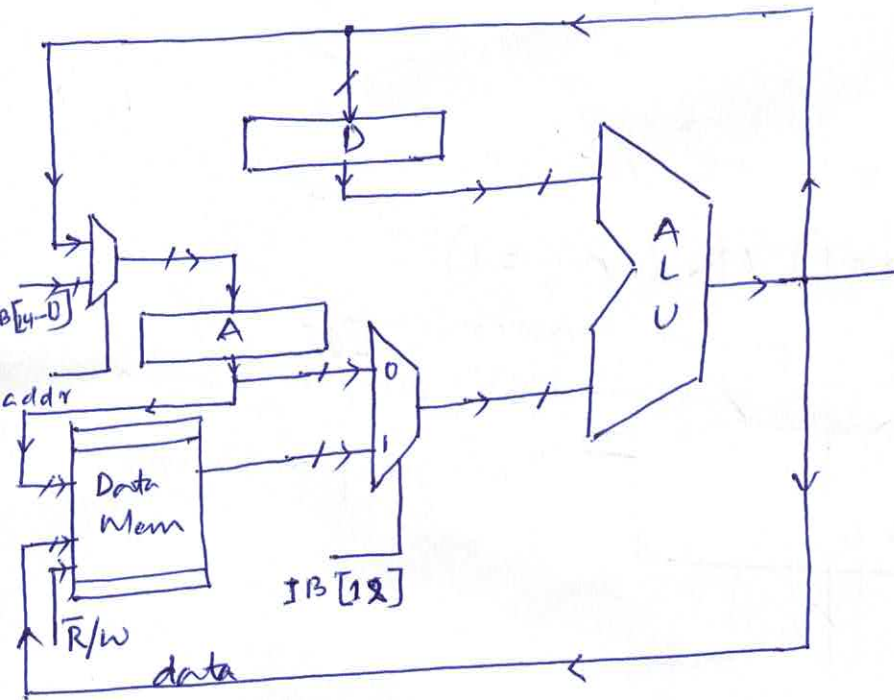
(8)

Writing back the result (WB) :-

- Writing back the result computed by ALU to respective destination.

Note:- Address/memory instruction doesn't require write back

ALU & Jump Instruction :-



updating program counter :-

